

# Hilbert Lattice Equations

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## OL, OML, modular OL

An *ortholattice* (OL) is an algebra  $\langle A, \vee, \wedge, ' \rangle$  in which the following conditions hold:

$$a \vee b = b \vee a \quad (1)$$

$$(a \vee b) \vee c = a \vee (b \vee c) \quad (2)$$

$$a'' = a \quad (3)$$

$$a \vee (a \wedge b) = a \quad (4)$$

$$a \wedge b = (a' \vee b')' \quad (5)$$

An *orthomodular lattice* (OML) is an OL in which the orthomodular law holds:

$$a \vee b = ((a \vee b) \wedge b') \vee b \quad (6)$$

A *modular OL* is an OL in which the modular law holds:

$$a \vee (b \wedge (a \vee c)) = (a \vee b) \wedge (a \vee c) \quad (7)$$

## $\mathcal{C}(\mathcal{H})$

A *Hilbert space* is a (for us, complex) vector space with an inner product, which is complete in the induced metric.

(See <http://us.metamath.org/mpegif/mmhil.html> for the complete list of axioms.)

The set of closed subspaces of a (possibly infinite-dimensional) Hilbert space  $\mathcal{H}$  is denoted  $\mathcal{C}(\mathcal{H})$ . It is a lattice; in particular, it is an orthomodular lattice (OML). This fact provides a primary motivation for studying the properties of OMLs.

$\mathcal{C}(\mathcal{H})$  is also called a *Hilbert lattice*.

## $\mathcal{C}(\mathcal{H})$ lattice operations

The orthocomplement  $a'$  of a closed subspace  $a$  (or actually any  $a \subseteq \mathcal{H}$ ) is the set of vectors orthogonal to all vectors in  $a$ :

$$a' \stackrel{\text{def}}{=} \{x \in \mathcal{H} : (\forall y \in a) (x, y) = 0\} \quad (8)$$

where  $(\cdot, \cdot)$  is the Hilbert vector space inner product. Note that  $a'$  is a closed subspace for any  $a \subseteq \mathcal{H}$ , and  $a''$  (the closure of  $a$ ) is the smallest closed subspace containing  $a$ .

The meet operation is just set-theoretical intersection:

$$a \wedge b \stackrel{\text{def}}{=} a \cap b \quad (9)$$

## $\mathcal{C}(\mathcal{H})$ lattice operations (cont.)

Ordering, join, unit, and zero can be defined in terms of these.  
 (We also define *commutes* and *Sasaki implication* for later use.)

$$a \leq b \stackrel{\text{def}}{\Leftrightarrow} a = a \wedge b \quad \Leftrightarrow \quad a \subseteq b \quad (10)$$

$$a \vee b \stackrel{\text{def}}{=} (a' \wedge b')' = (a + b)'' = (a \cup b)'' \quad (11)$$

$$0 \stackrel{\text{def}}{=} a \wedge a' = \{0\} = \mathcal{H}' \quad (12)$$

$$1 \stackrel{\text{def}}{=} 0' = \mathcal{H} \quad (13)$$

$$aCb \stackrel{\text{def}}{\Leftrightarrow} a = (a \wedge b) \vee (a \wedge b') \quad (\text{commutes}) \quad (14)$$

$$a \rightarrow_1 b \stackrel{\text{def}}{=} a' \vee (a \wedge b) \quad (\text{Sasaki implication}) \quad (15)$$

where  $+$  is subspace sum,  $\cup$  is set-theoretical union, and  $0$  is the zero vector. Note that  $\mathcal{C}(\mathcal{H})$  itself can be defined as

$$\mathcal{C}(\mathcal{H}) \stackrel{\text{def}}{=} \{x \subseteq \mathcal{H} : x = x''\} \quad (16)$$

## An open problem in $\mathcal{C}(\mathcal{H})$

“An open problem is to determine all equations satisfied by hilbertian lattices [i.e.  $\mathcal{C}(\mathcal{H})$ ], which would make possible the separation of the 'purely logic' part from the above axiomatics. It is not known if this problem is solvable, for it is not certain that these equations form a recursively enumerable set.”

—René Mayet, “Varieties of orthomodular lattices related to states,” *Algebra Universalis* 20 (1985), 368-396

## Equations known to hold in $\mathcal{C}(\mathcal{H})$

(See also last year's slide show)

- Orthomodular (OML) law (Husumi, 1937)
- Orthoarguesian (OA) law (Alan Day, 1975)
- Godowski's state-related equations (Godowski, 1981)
- Mayet's state-related equations (Mayet, 1985)
- $n$ -orthoarguesian ( $n$ -OA) laws (Megill/Pavičić, 2000)
- The modular law **does not hold** in ( $\infty$ -dimensional)  $\mathcal{C}(\mathcal{H})$

## Equations known to hold in $\mathcal{C}(\mathcal{H})$ (cont.)

OML law (Husumi, 1937):

$$((a \vee b) \wedge b') \vee b = a \vee b \quad (17)$$

OA law (Alan Day, 1975):

$$\begin{aligned} & a \wedge (((a \wedge b) \vee ((a \rightarrow_1 d) \wedge (b \rightarrow_1 d))) \vee \\ & \quad (((a \wedge c) \vee ((a \rightarrow_1 d) \wedge (c \rightarrow_1 d))) \wedge \\ & \quad ((b \wedge c) \vee ((b \rightarrow_1 d) \wedge (c \rightarrow_1 d)))) \leq b' \rightarrow_1 d \end{aligned} \quad (18)$$



## Equations known to hold in $\mathcal{C}(\mathcal{H})$ (cont.)

$n$ -OA example,  $n = 5$  (Megill/Pavičić, 2000):

$$\begin{aligned}
 & a \wedge (((((a \wedge b) \vee ((a \rightarrow_1 d) \wedge (b \rightarrow_1 d))) \vee \\
 & \quad (((a \wedge e) \vee ((a \rightarrow_1 d) \wedge (e \rightarrow_1 d))) \wedge \\
 & \quad ((b \wedge e) \vee ((b \rightarrow_1 d) \wedge (e \rightarrow_1 d)))))) \vee \\
 & \quad (((((a \wedge c) \vee ((a \rightarrow_1 d) \wedge (c \rightarrow_1 d))) \vee \\
 & \quad (((a \wedge e) \vee ((a \rightarrow_1 d) \wedge (e \rightarrow_1 d))) \wedge \\
 & \quad ((c \wedge e) \vee ((c \rightarrow_1 d) \wedge (e \rightarrow_1 d)))))) \wedge \\
 & \quad (((b \wedge c) \vee ((b \rightarrow_1 d) \wedge (c \rightarrow_1 d))) \vee \\
 & \quad (((b \wedge e) \vee ((b \rightarrow_1 d) \wedge (e \rightarrow_1 d))) \wedge \\
 & \quad ((c \wedge e) \vee ((c \rightarrow_1 d) \wedge (e \rightarrow_1 d)))))) \leq b' \rightarrow_1 d \quad (19)
 \end{aligned}$$

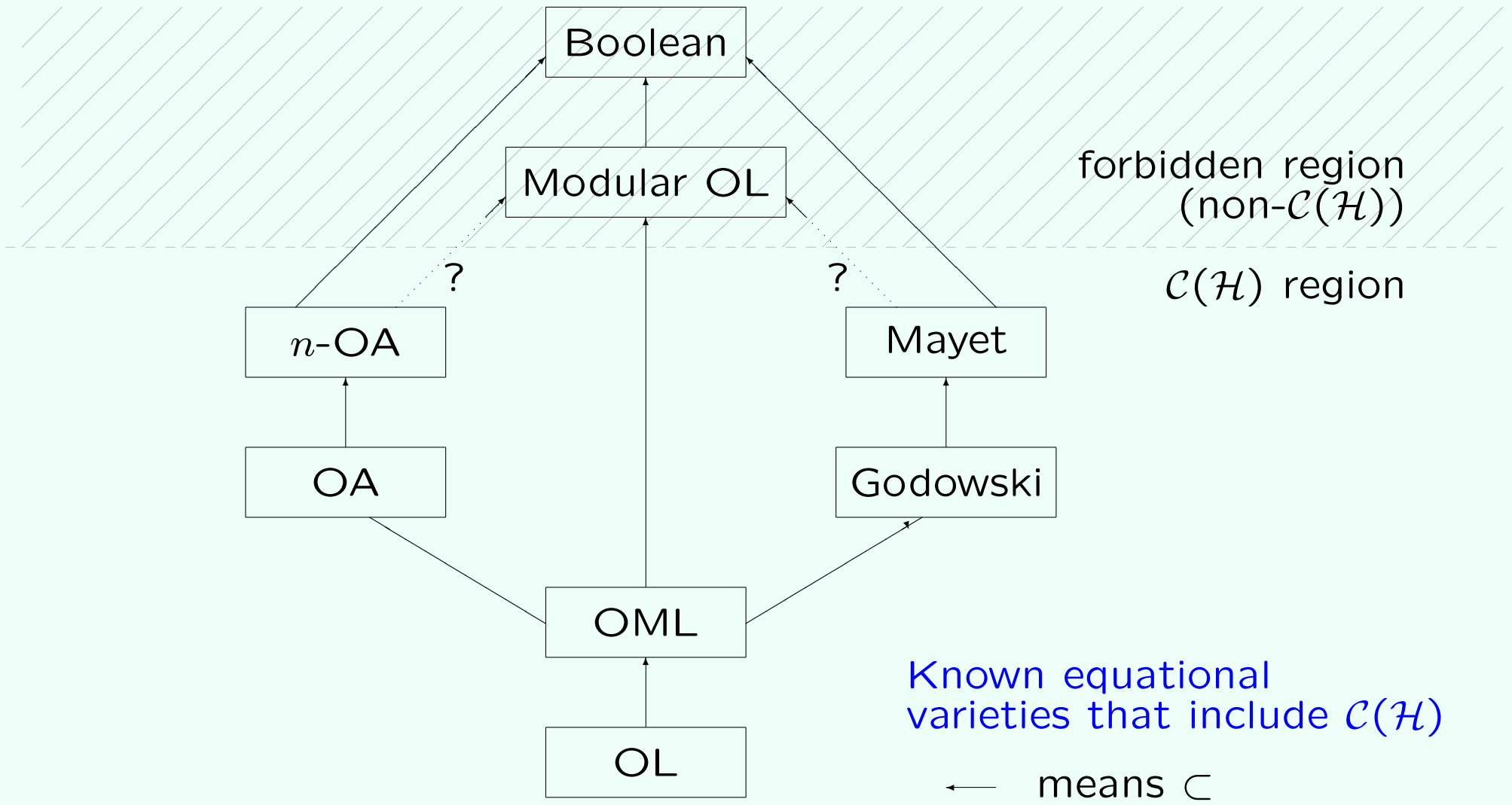
## Equations known to hold in $\mathcal{C}(\mathcal{H})$ (cont.)

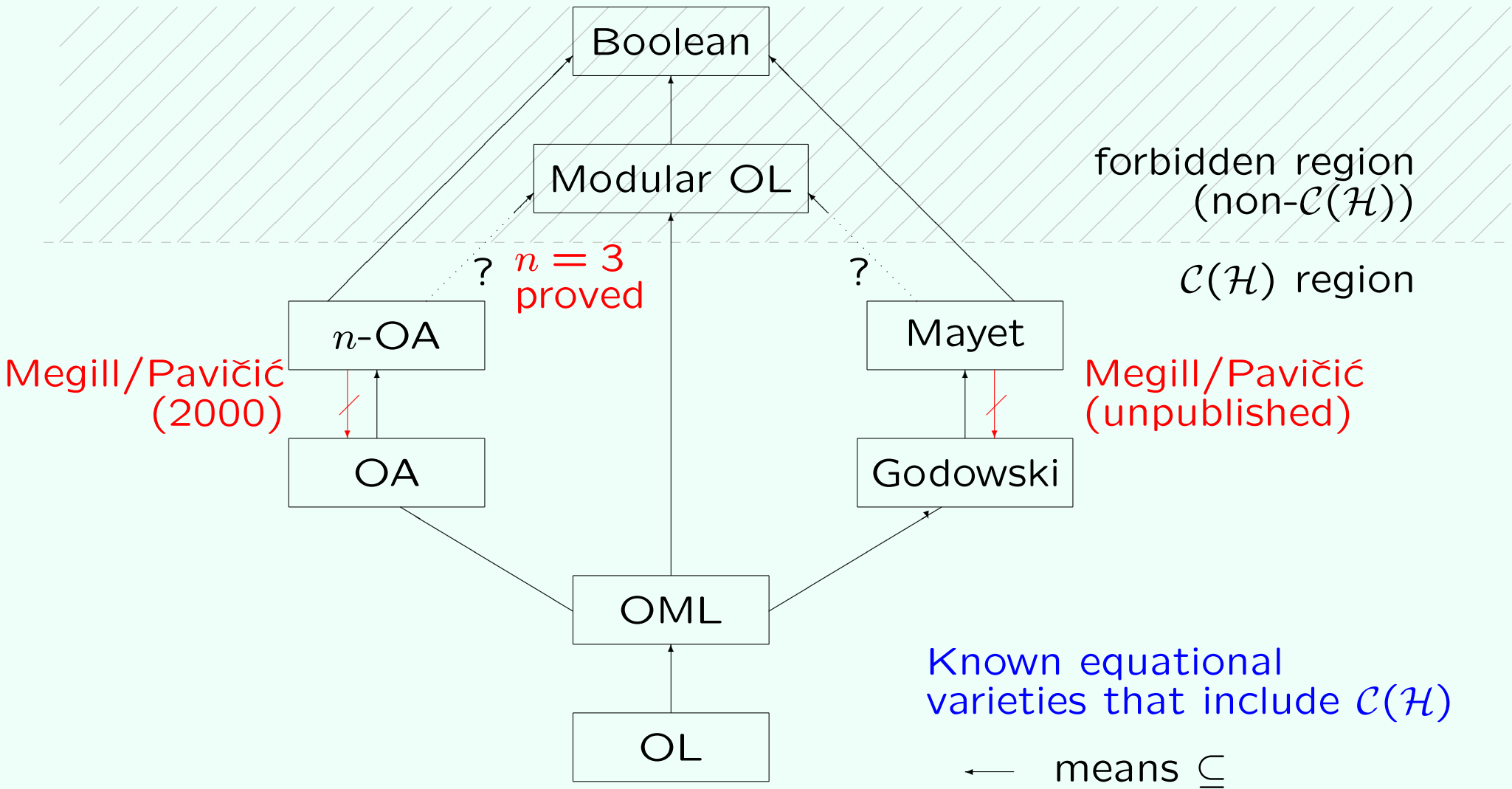
Example of Godowski equation (Godowski, 1981):

$$(a \rightarrow_1 b) \wedge (b \rightarrow_1 c) \wedge (c \rightarrow_1 a) \leq a \rightarrow_1 c \quad (20)$$

Example of Mayet equation (Megill/Pavičić, unpublished):

$$((a \rightarrow_1 b) \rightarrow_1 (c \rightarrow_1 b)) \wedge (a \rightarrow_1 c) \wedge (b \rightarrow_1 a) \leq c \rightarrow_1 a \quad (21)$$





## Modular pairs

The *modular pair* relation between two lattice elements  $a, b$ , denoted  $(a, b)M$ , is defined as

$$(a, b)M \stackrel{\text{def}}{\Leftrightarrow} (\forall x) [x \leq b \Rightarrow x \vee (a \wedge b) = (x \vee a) \wedge b] \quad (22)$$

The *dual modular pair* relation between two lattice elements  $a, b$ , denoted  $(a, b)M^*$ , is defined as

$$(a, b)M^* \stackrel{\text{def}}{\Leftrightarrow} (\forall x) [x \geq b \Rightarrow x \wedge (a \vee b) = (x \wedge a) \vee b] \quad (23)$$

## M-symmetry

The set of closed subspaces of infinite-dimensional Hilbert space,  $\mathcal{C}(\mathcal{H})$ , has a remarkable property: it is *M-symmetric*.

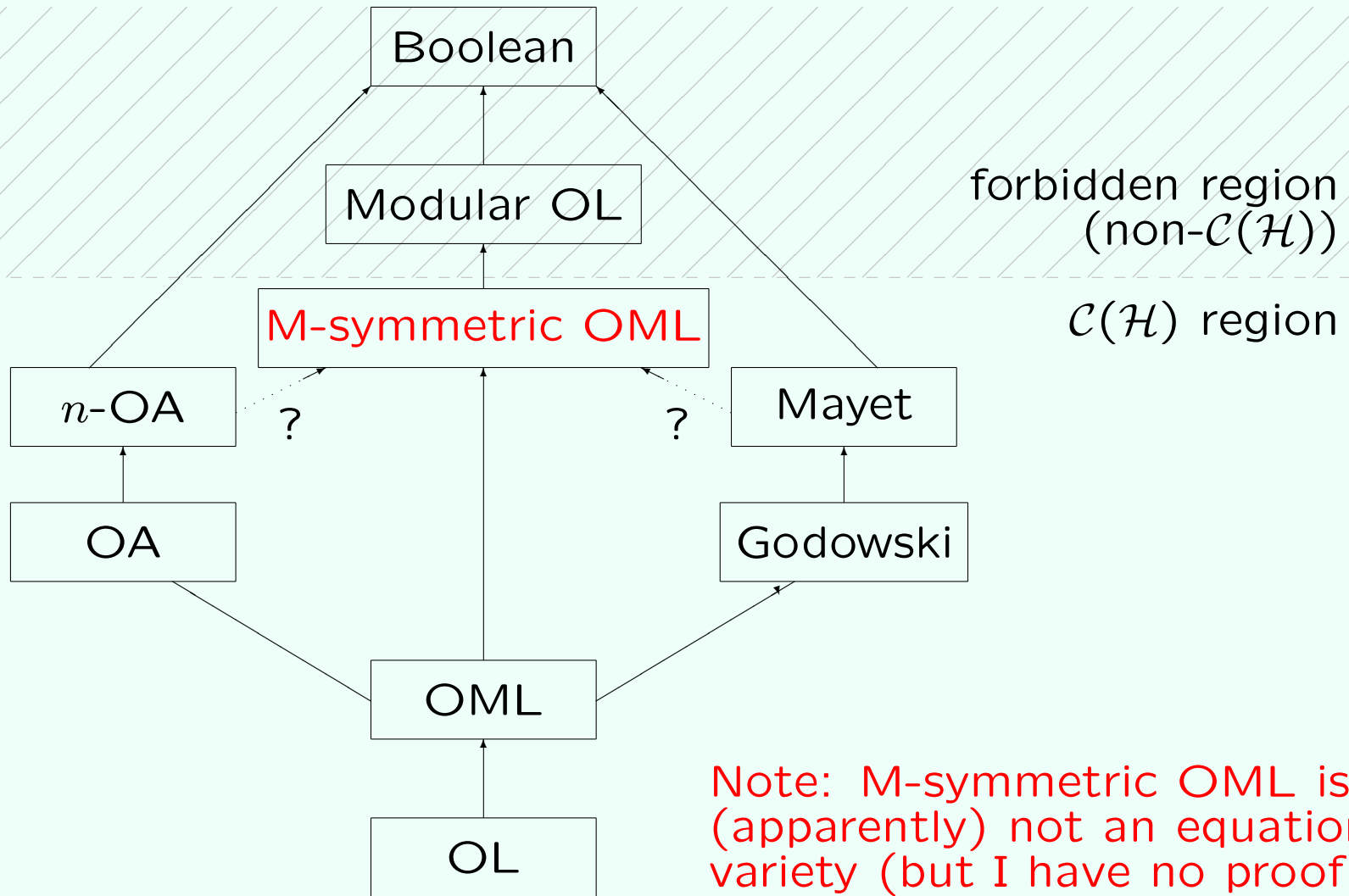
$$(a, b)M \Leftrightarrow (b, a)M \quad (24)$$

Other related symmetry properties also hold:

$$(a, b)M^* \Leftrightarrow (b, a)M^* \quad (M^*\text{-symmetric}) \quad (25)$$

$$(a, b)M \Leftrightarrow (b', a')M \quad (O\text{-symmetric}) \quad (26)$$

$$(a, b)M \Leftrightarrow (b, a)M^* \quad (\text{cross-symmetric}) \quad (27)$$



## Can M-symmetry help us find a new $\mathcal{C}(\mathcal{H})$ equation?

The M- (or M\*-)symmetry property is not an equation, because in prenex normal form it has an existential quantifier. To obtain an equation, one possible approach (that we are investigating) is to find a quantifier-free expression (polynomial equations connected with 'and')  $E(a, b, \dots)$  s.t.

$$E(a, b, \dots) \Rightarrow (b, a)M^* \quad (28)$$

holds in OML (or in some other known  $\mathcal{C}(\mathcal{H})$  condition). Then

$$E(a, b, \dots) \Rightarrow (a, b)M^* \quad (29)$$

will also hold in  $\mathcal{C}(\mathcal{H})$  and (after removal of  $(a, b)M^*$  quantifier) will be an equational inference that holds in  $\mathcal{C}(\mathcal{H})$ , hopefully stronger than the first condition.



## Our current working conjecture

We are trying to prove or disprove that the following equation holds in all OMLs:

$$\begin{aligned} p \wedge (a \vee (((a \wedge d) \vee (c \wedge (a \vee (c \wedge a'))))) \wedge d) \wedge b)) \\ \leq (((c \wedge d) \vee a) \vee p) \wedge b) \vee ((c \wedge d) \vee a) \end{aligned} \quad (30)$$

Expressed as a Mace4 problem:

```
% Use as the denial in the Mace4 example "nonmodular-oml.in"
(E ^ (A v (((A ^ D) v (C ^ (A v (C ^ c(A))))) ^ D) ^ B)))
v (((((C ^ D) v A) v E) ^ B) v ((C ^ D) v A))
!= (((C ^ D) v A) v E) ^ B) v ((C ^ D) v A).
```

## References

Most of the references for this material can be found at:

<http://us.metamath.org/qlegif/mmql.html#ref>

More miscellaneous stuff can be found at:

<http://us.metamath.org/award2003.html>